Progressive Iterative Interpolation for Subdivision Surfa
es

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Abstra
t. In this paper we give an insight to the interpolation method proposed by $[1]$. First, we point out that this method is an extension of the progressive iterative approximation of B-spline surfa
e to subdivision surface. We also solve the left open problem in $[1]$: proving the onvergen
e of the iterative interpolation for subdivision surfa
es. Then based on our analysis, we give a more simple and efficient modified interpolation algorithm for subdivision surfa
es.

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t Modeling - urve, surface, solid and object representations;

Keywords: subdivision surfa
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es, Catmull-Clark subdivision surfa
es, interpolation

Introduction 1

Subdivision surface is popular now in the fields of Computer Animation, CAD, Geometri Modeling, and so on. The ability to model arbitrary topology surfa
e makes it more suitable than classical spline surfaces in some applications. [3] proposed the Catmull-Clark subdivision surface which is the generalization of bicubic spline surface. [2] designed the Doo-Sabin subdivision method whi
h is the generalization of quadrati spline surface. Later, [9] developed the Loop subdivision for triangle mesh whi
h generalized the Box spline. All these three popular subdivision methods are approximating s
hemes. The other type of subdivision method is interpolating s
heme whi
h interpolates its original mesh. The famous interpolation subdivision method is the butterfly subdivision method proposed in [7] which was modified to generate smoother surfaces in [10]. The interpolating scheme for quadrilateral mesh was proposed in $[19]$.

Interpolation is one of general approaches in surface design and shape modeling. There are plenty of literatures dealing with the interpolation problem with different surface representations. As the appearance of re
ursive subdivision surfa
e, various interpolation methods based on subdivision surfa
es have been developed. One kind of those methods requires solving a

global linear system of equations, like $[5]$ and $[4]$. For the omputational ost of solving a large linear system of equations, many resear
hers developed other methods. [6] proposed an always worked method by using a two-phase subdivision method. The method of $[16]$ avoided the solving a linear equations uses the concept of similarity. And [8] gave one interpolation method avoiding solving a linear system of equations using the quasi-interpolation. Recently, $[1]$ presented a very simple and efficient interpolation approach by just moving the verti
es of the mesh. But they didn't provide the proof of its onvergen
e and left it as an open question

In this paper, we give insight to $[1]$'s method and give the convergence proof of this method. We find that this method ould be viewed as the extension of the progressive iteration interpolation for uniform spline proposed in $[14]$ and $[11]$. This is so called progressive iterative approximation property of uniform B-spline bases. More general results for non-uniform B-spline has been given in $[13]$. $[12]$ presented that if the given basis is totally positive, and its orresponding ollo
ation matrix is non-singular, this bases have the progressive iteration approximation property. [15] further proved that the normalized B-basis satisfies the progressive iterative approximation property with the fastest onvergen
e rates.

In this paper, we will prove that the subdivision surfa
es also satisfy the progressive iterative approximation property. In section 2, we will analyze the convergen
e for Loop subdivision surfa
e and give a modied algorithm. In se
tion 3, we present the similar result for Catmull-Clark subdivision surface restricted to quadrilateral mesh. Then we will give some examples to investigate the onvergen
e rate and give our on
lusion.

2 Iterative interpolation for Loop subdivision surfa
e:

The subdivision surfa
e for triangle mesh was proposed in [9]. It also analyzed its continuity and gave the formula to get the limit point on the limit surface. The limit points of a vertex with valence n as

Figure 1: A vertex with valence n in Loop subdivision

shown in Fig. 1 on the Loop subdivision surface can be al
ulated through the following formula:

$$
V_{\infty} = \beta_n V + (1 - \beta_n) Q \tag{1}
$$

where $\beta_n = \frac{3}{(1 - 2)(\beta + 3 + 1 - 2\pi)^2}$, $Q = \frac{1}{n} \sum_{i=1}^n Q_i$.

 $11-8\times$ $\frac{1}{8}$ + $\frac{1}{8}$ + $\frac{2}{4}$ cos $\frac{20}{n}$)² The essential part of the idea of progressive iterative interpolation is to find a corresponding point on the surfa
e for ea
h interpolated point and use the differences between them to update the control mesh. If the interpolated point is V and the corresponding point on the surface is V^k , the updating process could be written as:

$$
e^k = V - V^k
$$

$$
V^{k+1} = V^k + e^k
$$

Based on the above limit point formula, we can construct the iterative fitting algorithm for Loop subdivision surfa
e by using the limit point of ea
h vertex as the orresponding point. Assume the urrent mesh is M^k , and then we can get a mesh M^{k+1} which is a better approximation to original mesh. For ea
h vertex $V^{\,k}$ of $M^{\,k}$ point $V_{\infty}^{k} = \beta_{n} V^{k} + (1 - \beta_{n}) Q^{k}$. Then the difference is calculat as $e^{k} = V - V^{k}$. Now every vertex of M^{k+1} is updated by summing the corresponding vertex of M^k and its difference, that is $V^{k+1} = V^k + e^k$. As our below proof shows, the iteration will converge to a mesh which generates a Loop subdivision surface interpolating the original mesh. The details of the iterative interpolation algorithm are presented here:

Iterative Interpolation Algorithm

1. Input a triangle mesh M and the maximum error E_M ;

- 2. Initialize the $CurrentMaxError = E_M + 1$;
- 3. $k = 0$;
- 4. While $(CurrentMaxError > E_M)$ {

For every vertex V in M^{κ} {

Cal
ulating the orresponding vertex on the limit surface V^k_∞ ; $e^k = V - V^k_{\infty};$ $V^{k+1} = V^k + e^k$; if $(\|e^k\|)$ \leq $CurrentMaxError)$ $10ptCurrentMaxError = ||e^k||;$ \mathcal{F} 5. $k = k + 1$;

 $6.$ }

We can prove that e^{μ} always less than the e^{μ} , so it is not necessary to use the local adjusting in [1]. The method here is mu
h simpler than the algorithm proposed in [1], which needs a complicated shortest point omputation. The onvergen
e of our algorithm is analyzed in detail next.

2.1 Convergen
e of the iterative interpolation for Loop subdivision surface

Before our proof, a fa
t about the eigenvalues of the product of positive definite matrices needs to be mentioned. We present this fact in Lemma 1.

Lemma 1 If A and B are positive definite, the eigenvalues of AB are positive.

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The proof of Lemma 1 is quickly followed by the fact that if P and Q are square matrices of order n , then PQ and QP have the same eigenvalues (cf. Magnus and Neudecker [17], 1988, P.14).

We know that the differences of the $k + 1$ -th step e^{k+1} could be written as:

$$
e^{k+1} = V - V_{\infty}^{k+1}
$$

= $V - (\beta_n V k + 1 + (1 - \beta_n) Q^{k+1})$
= $V - (\beta_n V^{k+1} + (1 - \beta_n) Q^{k+1})$
 $- (\beta_n e^k + \frac{1 - \beta_n}{n} \sum_{i=1}^n e^k_{Q_i})$
= $e^k - (\beta_n e^k + \frac{1 - \beta_n}{n} \sum_{i=1}^n e^k_{Q_i})$

Thus in the matrix form, we have

$$
\begin{bmatrix} e_1^{k+1}, e_2^{k+1}, \dots, e_n^{k+1} \end{bmatrix}^T = (I - B) \begin{bmatrix} e_1^k \\ e_2^k \\ \vdots \\ e_n^k \end{bmatrix}
$$

$$
\begin{bmatrix} e_1^{k+1}, e_2^{k+1}, \dots, e_n^{k+1} \end{bmatrix}^T = (I - B)^k \begin{bmatrix} e_1^1 \\ e_2^1 \\ \vdots \\ e_n^1 \end{bmatrix}
$$

where I is one identity matrix and B is

$$
\begin{pmatrix}\n\beta_{n_1} & \dots & \frac{1-\beta_{n_1}}{n_1} & \dots \\
\vdots & \ddots & \vdots \\
\frac{1-\beta_{n_i}}{n_i} & \dots & \beta_{n_i} & \dots \\
\vdots & \vdots & \vdots \\
\end{pmatrix}
$$

The matrix B has the following properties:

(1)
$$
a_{ij} \ge 0
$$
, and $\sum_{j=1}^{n} a_{ij} = 1$, that is $||B||_{\infty} = 1$;

(2) There are n_i+1 positive elements in *i*-th row, and the positive elements in each row are equal except the element in the diagonal line.

(3) If
$$
a_{ij} = 0
$$
, then $a_{ji} = 0$

Properties (1) and (2) are simply implied from the formula of e^{k+1} , and property (3) is because the differences of two consecutive levels could be related to each other for their vertices constituting an edge in the mesh. Due to these properties, we can rewrite the matrix B as

$$
B=DS
$$

where D is diagonal matrix and S is a symmetric matrix

$$
\mathbf{D} = \begin{pmatrix} \frac{1-\beta_{n_1}}{n_1} & 0 & \cdots & 0 \\ 0 & \frac{1-\beta_{n_2}}{n_2} & \cdots & 0 \\ \vdots & & & \ddots & \\ 0 & & & & \frac{1-\beta_{n_k}}{n_k} \end{pmatrix}
$$

$$
\mathbf{S} = \begin{pmatrix} \frac{n_1\beta_{n_1}}{1-\beta_{n_1}} & \cdots & 1 & \cdots \\ \vdots & \ddots & & & \\ 1 & \cdots & \frac{n_i\beta_{n_i}}{1-\beta_{n_i}} & \cdots \\ \vdots & & & & \frac{n_k\beta_{n_k}}{1-\beta_{n_k}} \end{pmatrix}
$$

D is no problem positive definite for $\frac{1-\beta_{n_i}}{n_i} > 0$ ($1 \leq i \leq k$). Now we argue that the matrix *S* is also positive definite, which plays a key role in our convergence proof.

Proposition 1 The matrix S is positive definite.

Proof: To prove S is positive definite, we consider the corresponding quadric form.

$$
f(x_1, x_2, \ldots, x_n) = \mathbf{X}^T S \mathbf{X}
$$

where $X = (x_1, x_2, ..., x_n)^T$. If $f(x_1, x_2, ..., x_n)$ is positive for any none zero X , the symmetric matrix S is positive definite.

We notice that if vertices i and j are the endpoints of one edge e_{ij} in the mesh, then $a_{ij} = a_{ji} = 1$ in the matrix S . And an edge is adjacent to two faces exactly in a closed triangle mesh. Hence, we have

$$
f(x_1, x_2,..., x_n) = \sum_{\text{all faces}} \frac{1}{2} (x_i + x_j + x_r)^2 + \sum_{i=1}^k \left(\frac{n_i \beta_{n_i}}{1 - \beta_{n_i}} - \frac{n_i}{2} \right) x_i^2
$$

where x_i , x_j , and x_r are the corresponding three vertices of a face in the triangle mesh. From the formula 1, we know $\frac{n\beta_n}{1-\beta_n} \geq \frac{2}{3}n$ for $n \geq 3$. Hence, $f(x_1, x_2, \ldots, x_n)$ is positive for any none zero X. Thus S is positive definite.

Based on the above lemma and proposition, we can easily derive that the iterative interpolation for Loop subdivision is convergent.

Proposition 2 The iterative fitting algorithm for Loop subdivision surface is convergent.

Proof: The algorithm is convergent if and only if the absolute value of all eigenvalues of the matrix $P =$

ID: $papers_18$ Page: 3 Figure 2: A vertex with valence n in Catmull-Clark subdivision

 $I - B$ are less than 1. Hence, if all eigenvalues λ_i (1 < $i \leq k$) of B are $0 < \lambda_i \leq 1$, then all eigenvalues of P are $0 \leq 1 - \lambda_i < 1$ ($1 \leq i \leq k$). It means that the algorithm is onvergent.

Since $||B||_{\infty} = 1$, we have $\lambda_i \leq 1$. Also we have $B = DS$, where D and S are both positive definite. Hence all eigenvalues of B are positive. Thus, the iterative interpolation algorithm is onvergent.

3 Convergen
e of the iteration interpolation for Catmull-Clark Subdivision surfa
e

Based on the same idea, we can prove the convergence of iterative fitting algorithm for Catmull-Clark Subdivision surfa
e. Firstly, the orresponding limit point of a vertex with valen
e n in the mesh an be evaluated by the following equality.

$$
V_{\infty} = \frac{n^2}{n(n+5)}V + \frac{4}{n(n+5)}\sum_{i=1}^{n} E_i
$$

$$
+ \frac{1}{n(n+5)}\sum_{i=1}^{n} F_i
$$
(2)

where E_i is the edge point of vertex V, F_i is are the face points of vertex V . This evaluation formula requires that the fa
es surrounding vertex V are all quadrilateral fa
e. Be
ause our proof ompletely depends on this formula, our onvergen
e proof an only be applied on the quadrilateral mesh. It could be viewed as one deficiency of our proof. But in practi
e the quadrilateral mesh is also very popular, so our method has practical effects as well. Through the

formula 2, we can get the relation of the consecutive differences as following:

$$
e^{k+1} = V - V_{\infty}^{k+1}
$$

\n
$$
= V - \left(\frac{n^2}{n(n+5)}V^{k+1} + \frac{4}{n(n+5)}\sum_{i=1}^n E_i^{k+1}\right)
$$

\n
$$
+ \frac{1}{n(n+5)}\sum_{i=1}^n F_i^{k+1}\right)
$$

\n
$$
= V - \left(\frac{n^2}{n(n+5)}V^k + \frac{4}{n(n+5)}\sum_{i=1}^n E_i^k\right)
$$

\n
$$
+ \frac{1}{n(n+5)}\sum_{i=1}^n F_i^k - \left(\frac{n^2}{n(n+5)}e^k\right)
$$

\n
$$
+ \frac{4}{n(n+5)}\sum_{i=1}^n e_{E_i}^k + \frac{1}{n(n+5)}\sum_{i=1}^n e_{E_i}^k
$$

\n
$$
+ \frac{1}{n(n+5)}\sum_{i=1}^n e_{F_i}^k + \frac{4}{n(n+5)}\sum_{i=1}^n e_{E_i}^k
$$

\n
$$
+ \frac{1}{n(n+5)}\sum_{i=1}^n e_{F_i}^k
$$

Then using the matrix form, we have

$$
[e_1^{k+1}, e_2^{k+1}, \dots, e_n^{k+1}]^T = (I - C) \begin{bmatrix} e_1^k \\ e_2^k \\ \vdots \\ e_n^k \end{bmatrix}
$$

$$
[e_1^{k+1}, e_2^{k+1}, \dots, e_n^{k+1}]^T = (I - C)^k \begin{bmatrix} e_1^1 \\ e_2^1 \\ \vdots \\ e_n^1 \end{bmatrix}
$$

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Model	$\#$ of vertices	$#$ of iterations	Max Error	Ave Error
Horse	233	1 ດ	0.000912016	0.00022811
Bird	1129		0.000766811	8.8345e-5
Hand	6191		0.000914111	0.000164207
Boy	17342		0.000913795	9.5615e-5

Table 1: Interpolating results for Loop subdivision surfa
e.

where I is identity matrix and C is

$$
\begin{pmatrix}\n\frac{n_1^2}{n_1(n_1+5)} & \cdots & \frac{4}{n_1(n_1+5)} & \cdots & \frac{1}{n_1(n_1+5)} & \cdots \\
\vdots & \ddots & \vdots & \vdots & \vdots \\
\frac{4}{n_i(n_i+5)} & \frac{n_i^2}{n_i(n_i+5)} & \cdots & \frac{n_i^2}{n_i(n_i+5)} \\
\vdots & \vdots & \ddots & \vdots \\
\frac{1}{n_j(n_j+5)} & \cdots & \frac{n_j^2}{n_j(n_j+5)}\n\end{pmatrix}
$$

The matrix C could write as the product of two symmetric matrices.

$$
C=HM
$$

where H and M are

$$
\mathbf{H} = \begin{pmatrix} \frac{1}{n_1(n_1+5)} & 0 & \cdots & 0 \\ 0 & \frac{1}{n_2(n_2+5)} & \cdots & 0 \\ \vdots & & & & \\ 0 & & & & \frac{1}{n_k(n_k+5)} \end{pmatrix}
$$

$$
\mathbf{M} = \begin{pmatrix} n_1^2 & \cdots & 4 & \cdots & 1 & \cdots \\ \vdots & \ddots & & & \\ 4 & n_i^2 & & & \\ \vdots & & & & \\ 1 & & & & n_j^2 & \\ \vdots & & & & & \\ \end{pmatrix}
$$

For the matrix M , we know that if vertex v_i is the face point of v_i , then v_i is also one face point of v_i . This shows that if $a_{ij} = 1$, then $a_{ji} = 1$. Similarly, if $a_{ij} = 4$, then $a_{ji} = 4$ because one edge shares two endpoints. Obviously, we have $||C||_{\infty} = 1$. Thus the eigenvalues of C are not bigger than one. If we can prove that matrix M is positive definite, it is easy to show the iterative fitting algorithm for Catmull-Clark subdivision surfa
e is onvergent by following the same way for Loop subdivision surface. Actually, the matrix M is indeed positive definite. We give our proof here.

Proposition 3 The matrix M is positive definite.

Proof: The above observation shows why the matrix M is symmetri
. But these observations also give the hints for our proof. Using the definition of positive definite matrix, we write down the quadric form for M .

$$
f(x_1, x_2, \ldots, x_k) = X^T M X
$$

where $X = (x_1, x_2, \ldots, x_k)^T$. Thus we can rewrite $f(x_1, x_2, \ldots, x_k)$ as

$$
f = \sum_{\text{all faces}} (x_i + x_j + x_k + x_r)
$$

+
$$
\sum_{\text{all edges}} (x_i + x_j) + \sum_{\text{all vertices}} (n_i^2 - 3n_i) x_i^2
$$

We also know $n_i \geq 3$, hence $n_i^- - 3n_i \geq 0$. This means $f(x_1, x_2, \ldots, x_k) \geq 0$ for any none zero X. Thus, M is AT LEAST semipositive definite.

Following the same way of proving the convergence for Loop subdivision surface, we can easy know that the iterative interpolation algorithm for Catmull-Clark subdivision surface is also convergent. We conclude this in the following proposition.

Proposition 4 The iterative fitting algorithm for Catmul l-Clark subdivision surfa
e is onvergent.(NOT COMPLETELY CORRECT !)

$\overline{\mathbf{4}}$ **Results**

In this section, we give several examples. Table 1 shows the iteration numbers, maximum and average errors when using Loop subdivision surfa
es. As pointed out in [1], progressive iterative interpolation method is very efficient and can handle very huge meshes, and its complexity is only $O(mn)$ where n is the number of verti
es of the original mesh. Although no fairness ontrol fa
tor is added in progressive iterative interpolation, the results show that it ould produ
e visually pleasing surfa
e easily. So far the progressive iterative interpolation for Catmull-Clark subdivision surfa
e requires that the mesh is quadrilateral mesh. Ex
ept this limitation, the progressive iterative interpolation works very well. Figure 7 is the

result of this method applied on a mushroom mesh with 226 vertices. The maximum error for the mushroom mesh is 0.000813821, and the average error is $0.000235506.$

$\overline{5}$ Conclusion and future work

We have proved the convergence of iterative interpolation method for Loop and Catmull-Clark subdivision surfaces. This method could be treated as an extension of the progressive iterative approximation property from spline to subdivision. Based on our analysis, we give a modified progressive iterative interpolation algorithm. This modified algorithm is simpler and more efficient. But our proof to Catmull-Clark subdivision surface is only valid on the quadrilateral mesh. To extend our proof to arbitrary mesh for Catmull-Clark subdivision surface is our next concern. How to make this method effective for open mesh is also an interesting topic.

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 $\{a_i\}$ The original mesh with 233 vertices $\{a_i\}$ interpolating $\{a_i\}$ interpolating $\{a_i\}$

Figure 3: Loop subdivision surfa
e interpolating the horse mesh after 12 iterations.

(a) The original mesh with 1129 vertical contracts and the contract of the contract of the contracts of

Figure 4: Loop subdivision surfa
e interpolating the bird mesh after 9 iterations.

(a) The original mesh with ⁶¹⁹¹ verti
es (b) Interpolating Loop surfa
e

Figure 5: Loop subdivision surfa
e interpolating the hand mesh after 3 iterations.

(a) The original mesh with 17342 vertices and $\mathcal{C}^{(n)}$ interpolating Theorem (

Figure 6: Loop subdivision surfa
e interpolating the boy mesh after 6 iterations.

 $\{a\}$ The original mesh with 226 vertices $\{a\}$ is a construction of $\{a\}$ interpolating $\{a\}$ and $\{a\}$

Figure 7: Catmull-Clark subdivision surfa
e interpolating the mushroom mesh after 8 iterations.